



US007009901B2

(12) **United States Patent**
Baker

(10) **Patent No.:** **US 7,009,901 B2**
(45) **Date of Patent:** **Mar. 7, 2006**

(54) **SYSTEM AND METHOD FOR SENSING DATA STORED IN A RESISTIVE MEMORY ELEMENT USING ONE BIT OF A DIGITAL COUNT**

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(*) Notice: Subject to any disclaimer, the term of this patent is extended or adjusted under 35 U.S.C. 154(b) by 119 days.

(21) Appl. No.: **10/704,942**

(22) Filed: **Nov. 10, 2003**

(65) **Prior Publication Data**

US 2004/0095839 A1 May 20, 2004

Related U.S. Application Data

(62) Division of application No. 10/192,463, filed on Jul. 9, 2002, now Pat. No. 6,813,208.

(51) **Int. Cl.**
GIIC 7/02 (2006.01)

(52) **U.S. Cl.** **365/207**; 365/46; 365/148; 365/189.01; 365/236

(58) **Field of Classification Search** 365/207 O, 365/46 X, 148 X, 189.01 X, 236 X
See application file for complete search history.

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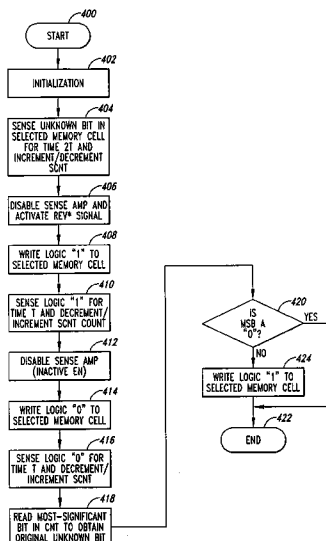
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(57) **ABSTRACT**

A method and system sense the logic state of an unknown initial data bit stored in a selected resistive memory cell. According to one method, a first count representing the logic state of the unknown initial data bit stored in the selected memory cell is generated. A second count is then generated, and represents a data bit having a first known logic state stored in the selected memory cell. A third count is then generated, and represents a data bit having a second known logic state stored in the selected memory cell. The logic state of the initial unknown data bit stored in the selected memory cell is then determined from the first, second, and third counts.

29 Claims, 5 Drawing Sheets



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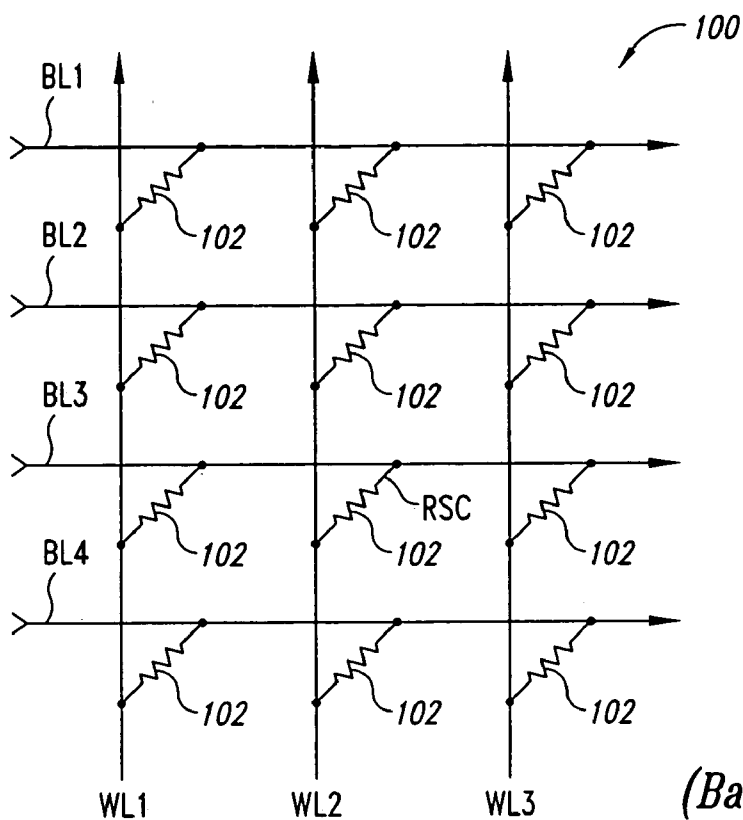


Fig. 1A
(Background Art)

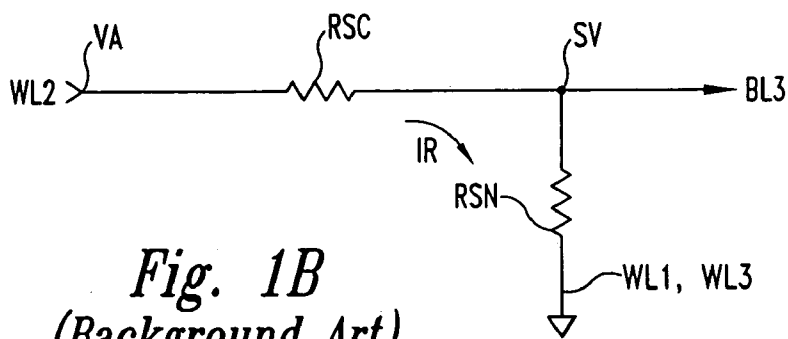


Fig. 1B
(Background Art)

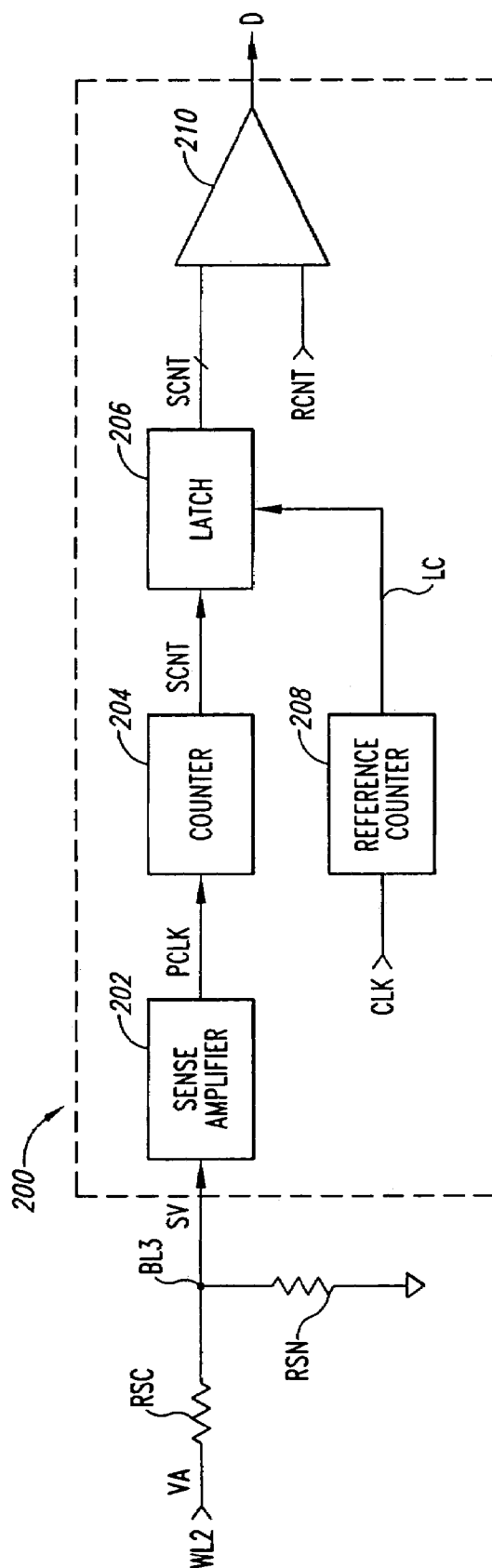


Fig. 2
(Background Art)

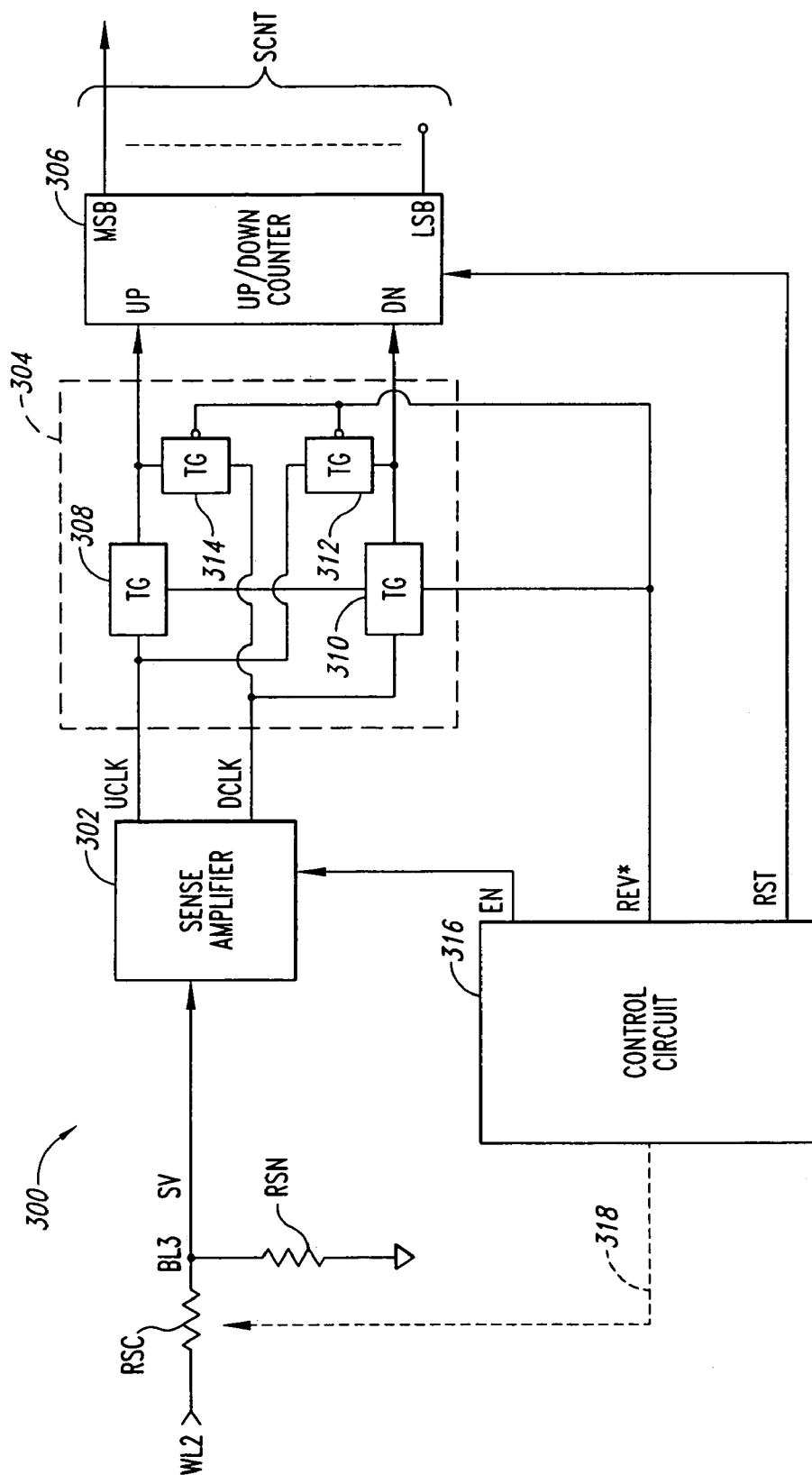
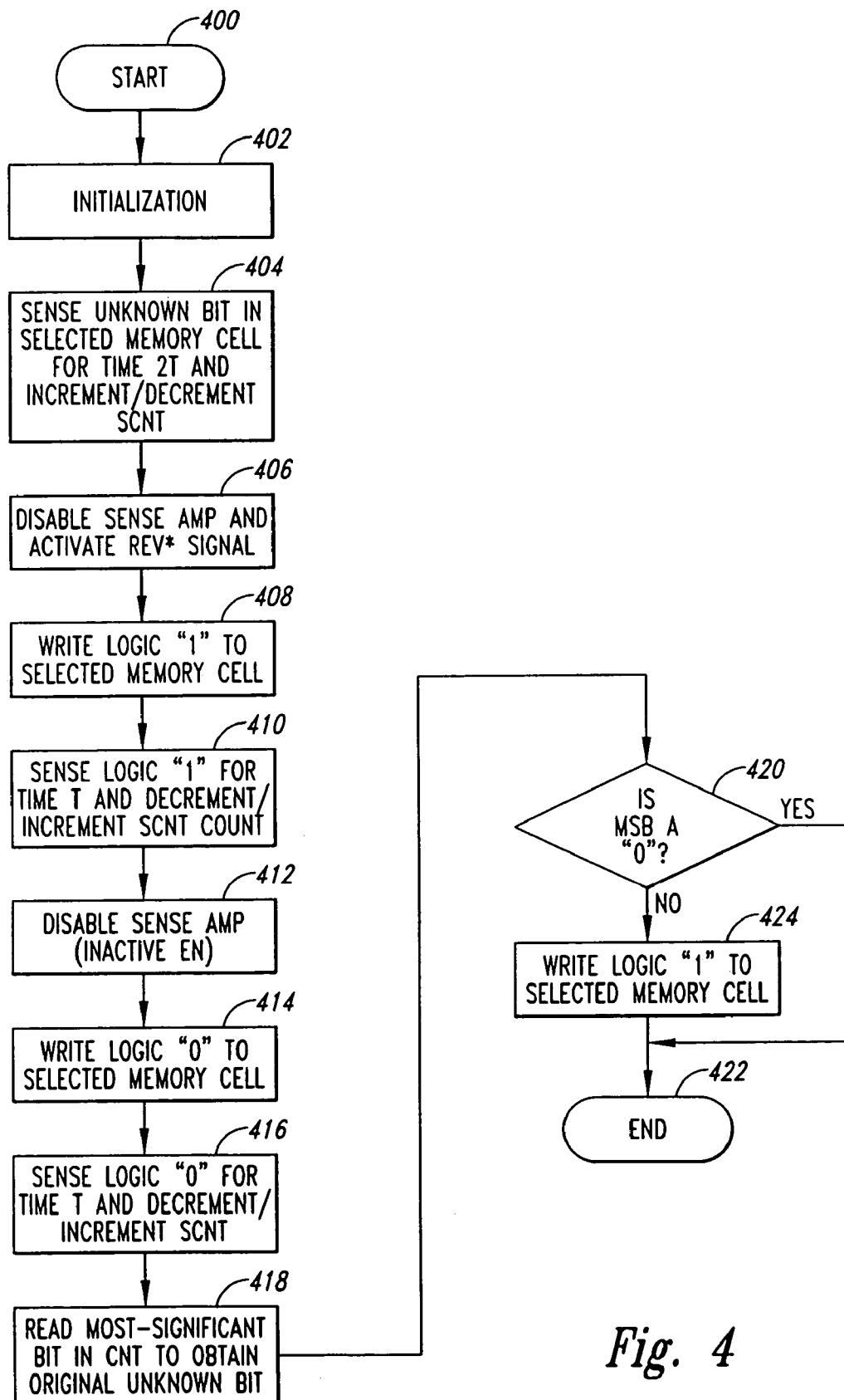
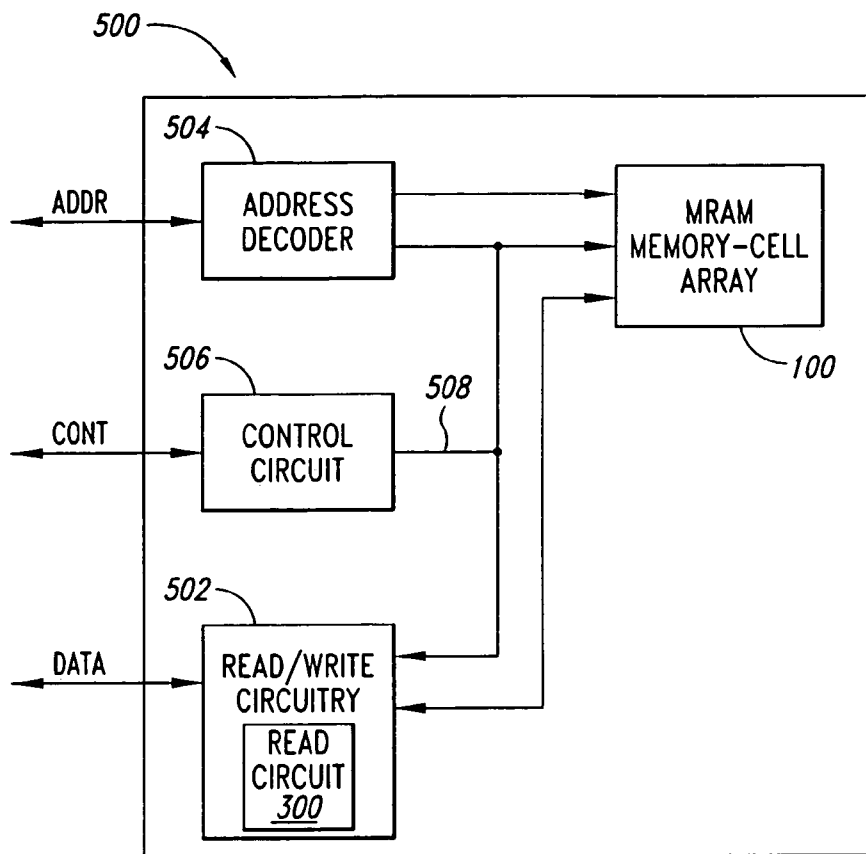
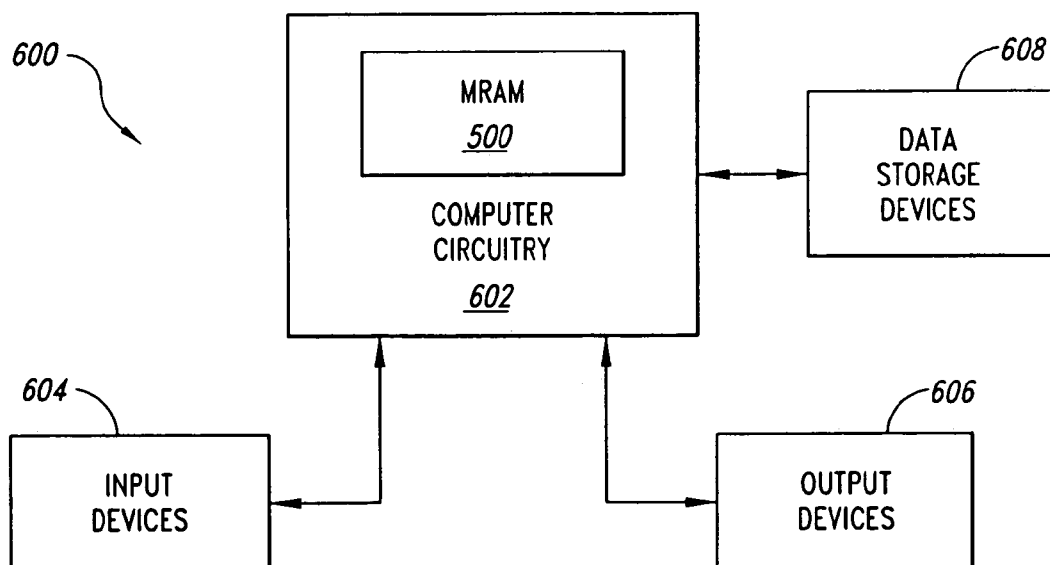


Fig. 3

*Fig. 4*

*Fig. 5**Fig. 6*

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SYSTEM AND METHOD FOR SENSING DATA STORED IN A RESISTIVE MEMORY ELEMENT USING ONE BIT OF A DIGITAL COUNT

CROSS-REFERENCE TO RELATED APPLICATION

This application is a divisional of U.S. patent application
Ser. No. 10/192,463, filed Jul. 9, 2002 now U.S. Pat. No. 6,813,208.

TECHNICAL FIELD

The present invention relates generally to integrated cir-
cuits, and more specifically to sensing data stored in an
integrated circuit memory device.

BACKGROUND OF THE INVENTION

Computer systems, video games, electronic appliances,
digital cameras, and myriad other electronic devices include
memory for storing data related to the use and operation of
the device. A variety of different memory types are utilized
in these devices, such as read only memory (ROM), dynamic
random access memory (DRAM), static random
access memory (SRAM), flash memory (FLASH), and mass
storage such as hard disks and CD-ROM or CD-RW drives.
Each memory type has characteristics that better suit that
type to particular applications. For example, DRAM is
slower than SRAM but is nonetheless utilized as system
memory in most computer systems because DRAM is
inexpensive and provides high density storage, thus allow-
ing large amounts of data to be stored relatively cheaply. A
memory characteristic that often times determines whether a
given type of memory is suitable for a given application is
the volatile nature of the storage. Both DRAM and SRAM
are volatile forms of data storage, which means the memo-
ries require power to retain the stored data. In contrast, mass
storage devices such as hard disks and CD drives are
nonvolatile storage devices, meaning the devices retain data
even when power is removed.

Current mass storage devices are relatively inexpensive
and high density, providing reliable long term data storage
relatively cheap. Such mass storage devices are, however,
physically large and contain numerous moving parts, which
reduces the reliability of the devices. Moreover, existing
mass storage devices are relatively slow, which slows the
operation of the computer system or other electronic device
containing the mass storage device. As a result, other
technologies are being developed to provide long term
nonvolatile data storage, and, ideally, such technologies
would also be fast and cheap enough for use in system
memory as well. The use of FLASH, which provides non-
volatile storage, is increasing popular in many electronic
devices such as digital cameras. While FLASH provides
nonvolatile storage, FLASH is too slow for use as system
memory and the use of FLASH for mass storage is imprac-
tical, due in part to the duration for which the FLASH can
reliably store data as well as limits on the number of times
data can be written to and read from FLASH.

Due to the nature of existing memory technologies, new
technologies are being developed to provide high density,
high speed, long term nonvolatile data storage. One such
technology that offers promise for both long term mass
storage and system memory applications is Magneto-Resis-
tive or Magnetic Random Access Memory (MRAM). FIG. 1

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is a functional diagram showing a portion of a conventional
MRAM array **100** including a plurality of memory cells **102**
arranged in rows and columns. Each memory cell **102** is
illustrated functionally as a resistor since the memory cell
has either a first or a second resistance depending on a
magnetic dipole orientation of the cell, as will be explained
in more detail below. Each memory cell **102** in a respective
row is coupled to a corresponding word line WL, and each
memory cell in a respective column is coupled to a corre-
sponding bit line BL. In FIG. 1, the word lines are desig-
nated WL1-3 and the bit lines designated BL1-4, and may
hereafter be referred to using either these specific designa-
tions or generally as word lines WL and bit lines BL. Each
of the memory cells **102** stores information magnetically in
the form of an orientation of a magnetic dipole of a material
forming the memory cell, with a first orientation of the
magnetic dipole corresponding to a logic "1" and a second
orientation of the magnetic dipole corresponding to a logic
"0." The orientation of the magnetic dipole of each memory
cell **102**, in turn, determines a resistance of the cell. Accord-
ingly, each memory cell **102** has a first resistance when the
magnetic dipole has the first orientation and a second
resistance when the magnetic dipole has the second orien-
tation. By sensing the resistance of each memory cell **102**,
the orientation of the magnetic dipole and thereby the logic
state of the data stored in the memory cell **102** can be
determined.

To write data to a selected memory cell **102**, a row current
IROW is applied to the word line WL coupled to the cell and
a column current ICOL is applied to the bit line BL coupled
to the cell. The row current IROW and column current ICOL
generated respective magnetic fields, with only the selected
memory cell **102** being subjected to both the magnetic field
generated by the row current and the magnetic field gener-
ated by the column current. The combination of these
magnetic fields applied to the selected memory cell **102** sets
the orientation of the magnetic dipole and thereby the
resistance of the cell, which writes a data bit corresponding
to either a logic 1 or 0 into the cell.

To read data from the MRAM array **100**, the resistance of
a selected memory cell **102** must be sensed. In one method
of sensing the resistance of a selected memory cell **102**, a
reference voltage VA is applied to the word line WL coupled
to the cell, and all other word lines and unselected bit lines
BL are coupled to ground. FIG. 1B is a schematic illustrating
the equivalent circuit of the MRAM array **100** when the
memory cell **102** coupled to the word line WL2 and bit line
BL3 is selected. In this situation, the reference voltage VA
is applied to the selected word line WL2, and all other word
lines WL1, WL3 and unselected bit lines BL1, BL2, BL4 are
coupled to ground. The resistance of the selected memory
cell **102** is represented by the resistance RSC, which is
coupled between word line WL2 and bit line BL3. All
unselected memory cells **102** coupled to the selected bit line
BL3 are coupled between the bit line BL3 and the unselected
word lines WL1, WL3, which are coupled to ground, and
these unselected memory cells collectively form a "sneak"
resistance RSN. All other unselected memory cells **102** in
the array **100** do not affect the equivalent circuit since both
ends of these memory cells are coupled to ground via the
unselected word lines WL1, WL3, and bit lines BL1, BL2,
and BL4, as will be appreciated by those skilled in the art.

In response to the applied reference voltage VA, a read
current IR flows through the resistance RSC presented by
the selected memory cell **102** and through the sneak resis-
tance RSN to ground. The current IR generates a sense
voltage SV on the selected bit line BL3, with the magnitude

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of this voltage being a function of the magnitude of the resistance RSC of the selected memory cell **102**. When the resistance RSC has a larger value, the sense voltage SV on the bit line BL3 will be less than when the resistance RSC has a smaller value. Accordingly, the sense voltage SV has a value indicating the magnitude of the resistance RSC and thus indicating the logic state of the data stored in the selected memory cell **102**.

In theory, sensing the resistance value of a selected memory cell **102** to read the logic state of data stored in the cell is simple as just described. In practice, however, reliable sensing is difficult due, in part, to the relatively small change in the resistance of the memory cell **102** between logic states. For example, in a typical MRAM array, each memory cell **102** has a resistance of about 1 Megaohm when the cell stores a logic "1" and a resistance of about 1.1 Megaohms when the cell stores a logic "0." The differential resistance of the selected memory cell **102** between a logic "1" and a logic "0" is thus only about 100 K Ω or approximately 10%. As a result, the sense voltage SV developed on the bit line BL3 varies by this same amount, making it difficult to reliably detect the sense voltage and determine whether a selected memory cell **102** stores a logic 1 or 0, as will be appreciated by those skilled in the art.

FIG. 2 is a functional block diagram illustrating an existing type of read circuit **200** for sensing the resistance values and thereby reading data from the memory cells **102** in the MRAM array **100** of FIG. 1. In FIG. 2, the same selected memory cell **102** discussed with reference to FIG. 1B is assumed to be selected and thus the equivalent circuit of FIG. 1B is shown as providing sense voltage SV as an input to the read circuit **200**. A sense amplifier **202** receives the sense voltage SV from the selected bit line BL3 and generates a pulse clocking signal PCLK which has a frequency of pulses that is a function of the value of the sense voltage SV. A counter **204** receives the PCLK signal and generates a sense count SCNT responsive to each pulse of the PCLK signal. The counter **204** applies the SCNT count to a latch **206** which latches the SCNT count responsive to a latch count signal LC from a reference counter **208**. The reference counter **208** is clocked by an applied clock signal CLK and activates the LC signal when an internal count equals a predetermined value, which occurs after a predetermined number of cycles of the CLK signal and thus after a predetermined time T. In this way, the reference counter **208** activates the LC signal to latch the SCNT every T seconds. The counter **204** resets the SCNT count responsive to the LC signal going active. A comparator **210** compares the latched SCNT from the latch **206** to a reference count RCNT and generates a data signal D responsive to this comparison. The RCNT count has a value corresponding to a threshold value for the resistance of the selected memory cell **102**. When the SCNT is greater than the RCNT count, the comparator **210** drives the data signal D high, and when the SCNT count is less than the RCNT count the comparator drives the data signal low.

In operation, the resistance RSC corresponding to the selected memory cell **102** and the sneak resistance RSN form a voltage divider, and develop the sense voltage SV on the selected bit line BL3 in response to the voltage VA applied to the selected word line WL2. The sense amplifier **202** generates the PCLK signal having a frequency determined by the value of the sense voltage SV. When the sense voltage SV has a first value corresponding to a first logic state, the PCLK signal has a first number of pulses over a given time period, and when the sense voltage has a second value corresponding to the complementary logic state, the

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PCLK signal has a second number of pulses over the time period. Initially, the SCNT count generated by the counter **204** the internal count of the reference counter **208** are set to 0 and the LC signal generated by the reference counter is inactive. In response to each pulse of the PCLK signal, the counter **204** increments the SCNT count, and at the same time the reference counter **208** increments the internal count in response to each rising-edge of the CLK signal.

The counter **204** continues incrementing the SCNT count in response to pulses of the PCLK signal. Because the PCLK signal has a frequency determined by the value of the sense voltage SV which, in turn, is determined by the resistance of the selected memory cell **102**, the SCNT count is incremented at a rate determined by the resistance of the selected memory cell. At the same time, the reference counter **208** increments the internal count responsive to the CLK signal. The counters **204**, **208** continue incrementing their respective counts in response to the PCLK and CLK signals until the internal count generated by the reference counter equals a predetermined value. The reference counter **208** increments the internal count once each cycle of the CLK signal, and thus the internal count equals a predetermined value after a predetermined number of cycles of the CLK signal, which occurs after a predetermined time T. Once the internal count of the reference counter **208** equals the predetermined value, the counter activates the LC signal causing the latch **206** to store the SCNT count at this point. The rate at which the SCNT count is incremented and thus the value of the latched SCNT count depends on the frequency of the PCLK signal which, in turn, depends on the value of the sense voltage SV. In this way, the counter **204** increments the SCNT count responsive to the PCLK signal for the time T, and at this point the value of the SCNT count is stored by the latch **206** and provided to the comparator **210**.

The comparator **210** compares the latched SCNT count to the reference count RCNT and drives the signal D either high or low depending on this comparison. When the latched SCNT count has a value that is less than the RCNT count, the sense voltage SV and thus the resistance RSC of the selected memory cell **102** corresponds to a first logic state and the comparator **210** drives the data signal D low, indicating the selected memory cell stores the first logic state. In contrast, when the latched SCNT count is greater than the RCNT count, the sense voltage and resistance RSC of the selected memory cell **102** correspond to the complementary logic state, and the comparator **210** drives the data signal D high, indicating the selected memory cell stores the complementary logic state.

In the read circuit **200**, the digital counts SCNT and RCNT are compared to ultimately determine the logic state of data stored in a selected memory cell **102**. As a result, the SCNT count must include enough bits and must be incremented at a sufficient rate responsive to the PCLK signal to ensure the generated RCNT has the desired resolution to enable the logic state of the selected memory cell **102** to be reliably sensed, as will be understood by those skilled in the art. Moreover, the circuitry required to form the counter **204** and comparator **210** will be relatively more complicated due to the need to generate and compare all bits in the RCNT and RCNT counts. In addition, with the read circuit **200** the SCNT count corresponding to a give logic state may vary among cells **102** in the array **100** (FIG. 1) due to variations in electrical characteristics among the cells resulting from the physical construction of the array. As a result, the RCNT may ideally need to be varied depending on which cell **102** is being selected, which further complicates the circuitry of the read circuit **200**. Although the above description focuses

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on MRAM technology, the concepts and principles are equally applicable to other types of resistive memory technologies.

There is a need for a simplified method and system for sensing the resistance value of resistive memory cells such as MRAM cells to reliably read data from the cells.

SUMMARY OF THE INVENTION

According to one aspect of the present invention, a method of sensing data stored in a selected resistive memory cell includes generating a first count representing the logic state of an unknown initial data bit stored in the selected memory cell. A second count is generated and represents a data bit having a first known logic state stored in the selected memory cell. A third count is generated and represents a data bit having a second known logic state stored in the selected memory cell. The logic state of the initial unknown data bit stored in the selected memory cell is then determined from the first, second, and third counts.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1A is a functional diagram showing a portion of a conventional MRAM array including a plurality of memory cells arranged in rows and columns.

FIG. 1B is a schematic illustrating the equivalent circuit for the MRAM array of FIG. 1A when a particular memory cell is selected during a read operation.

FIG. 2 is a functional block diagram illustrating a known type of read circuit for sensing the resistance values and thereby reading data from the memory cells in the MRAM array of FIG. 1.

FIG. 3 is a functional block diagram illustrating a read circuit for sensing the resistance values and thereby reading data from the memory cells in the MRAM array of FIG. 1 according to one embodiment of the present invention.

FIG. 4 is flowchart illustrating the overall process executed by the read circuit of FIG. 3.

FIG. 5 is a functional block diagram of an MRAM including the read circuit of FIG. 3.

FIG. 6 is a functional block diagram of a computer system including computer circuitry containing the MRAM of FIG. 5.

DETAILED DESCRIPTION OF THE INVENTION

FIG. 3 is a functional block diagram illustrating a read circuit 300 that generates a sense count SCNT having a most significant bit (MSB) corresponding to the logic state of the data stored in a selected memory cell 102 in the MRAM array 100 of FIG. 1 according to one embodiment of the present invention. In contrast to the prior data circuit 200 of FIG. 2, which generates an N-bit count SCNT that must then be compared to another N-bit reference count RCNT to determine the logic state of the data stored in a selected memory cell 102, the read circuit 300 generates single N-bit count SCNT having a single MSB bit indicating the logic state of data stored in the selected memory cell, as will be explained in more detail below. By eliminating the need to utilize the reference count RCNT and the need to compare two N-bit counts, the data circuit 300 provides a simplified and reliable circuit for reading data from MRAM memory cells. In the following description, certain details are set forth to provide a sufficient understanding of the present invention. However, it will be clear to one skilled in the art

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that the present invention may be practiced without these particular details. In other instances, well-known circuits and their operation have not been shown or described in detail to avoid unnecessarily obscuring the present invention.

In the read circuit 300, the same memory cell 102 discussed with reference to FIG. 1B is once again assumed to be selected, and thus the equivalent circuit of FIG. 1B is shown as providing the sense voltage SV as an input to the read circuit. A sense amplifier 302 receives the sense voltage SV from the selected bit line BL3 and generates an up clocking signal UCLK and a down clocking signal DCLK in response to the sense voltage. The sense amplifier 302 pulses either the UCLK or DCLK signal depending on the value of the sense voltage SV. When the sense voltage SV has a value greater than a threshold voltage, the sense amplifier 302 pulses the UCLK signal, and, conversely, when the sense voltage is less than the threshold voltage the sense amplifier pulses the DCLK signal. The sense amplifier 302 operates in this manner when an applied enable signal EN is active, and terminates generation of the UCLK and DCLK signals regardless of the applied sense voltage SV when the enable signal is inactive.

The value of the sense voltage SV varies around the threshold voltage due to inherent noise on the bit line BL3 and the small magnitude of the sense voltage, as will be appreciated by those skilled in the art. Thus, the sense amplifier 302 generates a series of UCLK, DCLK signal pulses, with the number of pulses of the UCLK signal and the number of pulses of the DCLK signal over a sensing time T being determined by the value of the sense voltage SV. When the sense voltage SV is greater than the threshold voltage, the sense amplifier 302 generates UCLK signal pulses and when the sense voltage is less than the threshold voltage the sense amplifier generates DCLK signal pulses. In the embodiment of FIG. 3, the resistance RSC of the selected memory cell 102 is assumed to have a value of approximately 1 Megaohm when the cell stores a logic 1, and a value of approximately 1.1 Megaohms when the cell stores a logic 0. As a result, the sense voltage SV has a larger value when the selected memory cell 102 stores a logic 1 and when the cell stores a logic 0. Thus, when the selected memory cell 102 stores a logic 1 the sense amplifier 302 will generate more UCLK signal pulses than DCLK signal pulses over the sensing time T.

The UCLK, DCLK signal pulses are applied through a switching circuit 304 to clock and up/down counter 306 which generates an N-bit sense count SCNT in response to the applied signal pulses. The SCNT count of the counter 306 is initially set to an initial value IV, which is thereafter incremented in response to UCLK signal pulses applied to an up input UP of the counter, and is decremented in response to DCLK signal pulses applied to a down input DN of the counter, as will be described in more detail below. The switching circuit 304 includes a plurality of transmission gates 308-314 coupled as shown between the sense amplifier 302 and the up/down counter 306. The transmission gates 308-314 operate in response to a reverse signal REV* from a control circuit 316 to either apply the UCLK and DCLK signals to the UP and DN inputs, respectively, of the counter 306, or to reverse the UCLK and DCLK signals and apply the UCLK signal to the DN input and the DCLK signal to the UP input. When the REV* signal is inactive high, the transmission gates 312, 314 turn OFF and the transmission gates 308, 310 turn ON, applying the UCLK signal to the UP input and the DCLK signal to the DN input, respectively, of the counter 306. In contrast, when the REV* signal is active

low, the transmission gates **308,310** turn OFF and the transmission gates **312, 314** turn ON, applying the UCLK signal to the DN input and the DCLK signal to the UP input, respectively, of the counter **306**.

In addition to the REV* signal, the control circuit **316** also generates a plurality of control signals **318** to control the logic state of data stored in the selected memory cell **102** represented by the resistance RSC, as will be explained in more detail below. The control circuit **316** also generates a reset signal RST that is applied to the counter **306** to set the SCNT count to the initial value IV and an enable signal EN is applied to the sense amplifier **302** to enable operation of the sense amplifier. One skilled in the art will understand circuitry for forming the sense amplifier **302**, transmission gates **308–314**, counter **306**, and control circuit **316**, and thus, for the sake of brevity, these components will not be described in more detail.

In operation, the read circuit **300** operates under control of the control circuit **316** to sense the logic state of data stored in the selected memory cell **102** represented by the resistance RSC, as will now be explained in more detail with reference to FIG. 3 and the flowchart of FIG. 4, which illustrates the process executed by the read circuit **300** in sensing the data stored in the selected memory cell. The process begins in step **400** and proceeds immediately step **402**, in which the control circuit **316** activates the RST signal to cause the counter **306** to reset the SCNT count to the initial value IV. The initial value IV is defined by the most significant bit MSB of the SCNT count being a 1 and all other bits being a 0. For example, if the SCNT count is 10 bits, the initial value IV equals “1000000000” with the leftmost “1” being the MSB. The control circuit **316** also activates the EN signal to enable the sense amplifier **302**, and drives the REV* signal inactive high to turn ON transmission gates **308, 310** and apply the UCLK, DCLK signals to the UP, DN inputs of the counter **306**.

From step **402**, the process proceeds to step **404** and the control circuit **316** accesses the selected memory cell **102** corresponding to the resistance RSC. In accessing the selected memory cell **102**, the control circuit **316** controls the MRAM array **100** (FIG. 1) as previously described to develop the sense voltage SV on the bit line BL3. As previously discussed, the sense voltage SV has a value determined by the resistance RSC, which corresponds to the data stored in the selected memory cell **102**. Note that at this point, it is unknown whether the selected memory cell **102** represented by the resistance RSC stores a logic 1 or logic 0. The sense amplifier **302** receives the sense voltage SV and pulses either the UCLK or DCLK signal depending on the value of the sense voltage SV. As previously described, if the sense voltage SV is greater than a threshold voltage the sense amplifier **302** pulses the UCLK signal, and if the sense voltage is less than the threshold voltage the sense amplifier pulses the DCLK signal. Each UCLK signal pulse is applied through the transmission gate **308** to the UP input of the counter **306**, incrementing the SCNT count responsive to each such pulse. Each DCLK signal pulse is applied through the transmission gate **3108** to the DN input of the counter **306**, which decrements the SCNT count responsive to each such pulse.

The read circuit **300** continues operating in step **404** for the time 2T, at which time the SCNT count generated by the counter **306** has a value that differs from the initial value IV by the net of the number of UCLK signal pulses received and the number of DCLK signal pulses receive. For example, when a logic 1 is stored in the selected memory cell **102**, the sense amplifier **302** may generate 251 UCLK

signal pulses and 249 DCLK signal pulses during each time T. Thus, after each time T, the net of the UCLK and DCLK signal pulses causes the counter **306** to increment the SCNT count by 2 (251 up pulses minus 249 down pulses equals a net of 2 up pulses). In this situation, after the time 2T the counter **306** as incremented the initial value IV of the SCNT count by 4 to IV+4. To illustrate another example, assume that when a logic 0 is stored in the selected memory cell **102**, the sense amplifier **302** generates 245 UCLK signal pulses and 255 DCLK signal pulses each time T. In this situation, after each time T the counter **306** decrements the SCNT count by 10 (245 up pulses minus 255 down pulses equals a net of 10 down pulses), and thus after the time 2T the counter as decremented the SCNT count from the initial value IV to the value IV–20. Note that in this example, the number of UCLK signal pulses is greater when the selected memory cell **102** stores a logic 1 than when the cell stores a logic 0 because the sense voltage SV is greater in the former situation, as previously discussed.

After the time 2T, the process proceeds to step **406** and the control circuit **316** drives the EN signal inactive, disabling the sense amplifier **302** so that the counter **306** maintains the SCNT count at the value generated after the time 2T, which was IV+4 in the case of a logic 1 being sensed in the example set forth above. At this point, the control circuit **316** also drives the REV* signal active low, turning OFF transmission gates **308, 310** and turning ON transmission gates **312, 314**. As a result, the UCLK signal pulses are now applied to the DN input of the counter **306** and the DCLK signal pulses applied to the IN input. From step **406** the process proceeds to step **408** and the control circuit **316** generates the control signals **318** to write a logic 1 into the selected memory cell **102** represented by the resistance RSC. As previously mentioned, the resistance RSC has a value of approximately 1 Megahom when storing a logic 1.

After storing a logic 1 in the selected memory cell **102**, the process goes to step **410** and the control circuit **316** activates the EN signal to enable the sense amplifier **302**. The sense amplifier **302** thereafter operates as previously described to generate the UCLK and DCLK signal pulses in response to the sense voltage SV corresponding to the logic 1 now stored in the selected memory cell **102**. Note that the active REV* signal causes the switching circuit **304** to apply the UCLK and DCLK signal pulses to the DN and UP inputs, respectively, of the counter **306**, which is the reverse of the prior situation during the initial sensing time 2T. The control circuit **316** maintains the EN signal active for the time T, and thus the sense amplifier **302** detects the sense voltage SV corresponding to the known logic 1 in the selected memory cell **102** and generates the corresponding UCLK and DCLK signal pulses for the time T.

Because the UCLK and DCLK signal pulses are applied to opposite inputs DN, UP of the counter **306**, the counter decrements the value of the SCNT count in response to each UCLK signal pulse and increments the value of the SCNT count in response to each DCLK pulse. As a result, the counter **306** adjusts the value of the SCNT count after the time T to remove a portion of the net count that would have been included in the SCNT count if a logic 1 was being sensed during the initial time 2T. In the example described above, after each time T the SCNT count is incremented by 2 if a logic 1 was stored in the selected memory cell **102** to a value IV+4 after the time 2T. Thus, after the time T with the counter **306** operating in the reverse direction, the counter **306** will have decremented the SCNT count by 2. The counter **306** thus decrements the SCNT count by 2 after each time T instead of incrementing the count by 2 due to the

UCLK and DCLK signal pulses being applied to opposite inputs DN, UP of the counter. Note that if the unknown data bit stored in the selected memory cell **102** during the initial time $2T$ was a logic 0, the counter **306** nonetheless decrements the SCNT count by 2 after the time T with the counter operating in the reverse direction.

After the sense amplifier **302** has sensed the known logic 1 for the time T , the process goes to step **412** and the control circuit **316** drives the EN signal inactive, disabling the sense amplifier **302** so that the counter **306** maintains the SCNT count at the value generated at the end of the last time T . At this point, the control circuit **316** maintains the REV* signal active low. From step **412** the process proceeds to step **414** and the control circuit **316** generates the control signals **318** to write a logic 0 into the selected memory cell **102** represented by the resistance RSC. As previously mentioned, the resistance RSC has a value of approximately 1.1 Megahoms when storing a logic 0. After storing a logic 0 in the selected memory cell **102**, the process goes to step **416** and the control circuit **316** once again activates the EN signal to enable the sense amplifier **302**. The sense amplifier **302** thereafter operates as previously described to generate the UCLK and DCLK signal pulses in response to the sense voltage SV corresponding to the logic 0 now stored in the selected memory cell **102**. Note that the active REV* signal once again causes the switching circuit **304** to apply the UCLK and DCLK signal pulses to the DN and UP inputs, respectively, of the counter **306**, which once again cause the counter **306** to operate the reverse of the operation during the initial sensing time $2T$. The control circuit **316** once maintains the EN signal active for the time T , and thus the sense amplifier **302** detects the sense voltage SV corresponding to the known logic 0 in the selected memory cell **102** and generates the corresponding UCLK and DCLK signal pulses for the time T .

Because the UCLK and DCLK signal pulses are once again applied to opposite inputs DN, UP of the counter **306**, the counter adjusts the value of the SCNT count to remove a portion of the net count that would have been included in the SCNT count if a logic 0 was being sensed during the initial time $2T$. In the example described above, after each time T the SCNT count is decremented by 10 if a logic 0 was stored in the selected memory cell **102**. Thus, after the time T with the counter **306** operating in the reverse direction, the counter **306** will have incremented the SCNT count by 10. The counter **306** thus increments the SCNT count by 10 after each time T instead of decrementing the count by 10 due to the UCLK and DCLK signal pulses being applied to opposite inputs DN, UP of the counter.

After the sense amplifier **302** has sensed the known logic 0 for the time T , the process goes to step **418** and the control circuit **316** drives the EN signal inactive, disabling the sense amplifier **302** so that the counter **306** maintains the SCNT count at the value generated at the end of the last time T . At this point, the most significant bit MSB of the SCNT count corresponds to a logic state of the data initially stored in the selected memory cell **102**. Thus, if the most significant bit MSB is a logic 1, the selected memory cell **102** initially stored in a logic 1. Conversely, if the most significant bit MSB is a logic 0, the selected memory cell **102** initially stored a logic 0. The counter **306** provides the most significant bit MSB as the output of the read circuit **300**. From step **418**, the process goes to step **420** and the control circuit **316** determines whether the most significant bit MSB of the SCNT count equals a logic 0, which indicates the selected memory cell **102** initially stored a logic 0. Note that at this point the selected memory cell **102** stores a logic 0 because

the control circuit **316** wrote a logic 0 to the selected memory cell in step **414**. As a result, if the selected memory cell **102** initially stored a logic 0, the process proceeds immediately to step **422** and terminates since the selected memory cell properly stores a logic 0. In contrast, if the control circuit **316** determines the most significant bit MSB of the SCNT count equals a logic 1, meaning the selected memory cell initially stored a logic 1, the process goes to step **424** and the control circuit rights a logic 1 to the selected memory cell to restore the cell to its proper initial logic state. From step **424**, the process then goes to step **422** and terminates.

The read circuit **300** senses an unknown bit stored in a selected memory cell **102** for a time $2T$ and generates a corresponding SCNT count. The SCNT count generated after the time $2T$ will have a unique value depending on whether the selected memory cell **102** stores a logic 1 or a logic 0. After the time $2T$, the value of the SCNT count is twice what the value would have been after the time T , with the precise value of the SCNT count being determined by whether the selected memory cell **102** stores a logic 1 or 0. After having generated this value of the SCNT count corresponding to the unknown bit stored in the selected memory cell **102**, the read circuit **300** subtracts count values corresponding to a known logic 1 and a known logic 0 stored in the selected memory cell. By subtracting the count values for a known logic 1 and 0, and knowing that a logic 1 stored in a selected memory cell **102** generates more UCLK signal pulses than a logic 0, the read data circuit generates the SCNT count having the MSB bit indicating the logic state of data stored in the selected memory cell **102**. To provide a better understanding of the operation of the read circuit **300**, an example will now be discussed with reference to Table 1 set forth below.

TABLE 1

(Logic 1 = 251 UCLK, 249 DCLK; Logic 0 = 245 UCLK, 255 DCLK)			
Time	Sensed Bit	SCNT Count	Decimal Value
Start		1000000000	512
T	Unknown "1"	1000000010	514
T	Unknown "1"	1000000100	516
Switch UP and DN			
T	Known "1"	1000000010	514
T	Known "0"	1000001100	524

Table 1 illustrates the SCNT counts generated by counter **306** in the read circuit **300** after respective times T whether a logic 1 and 0 stored in the selected memory cell **102** have the specified characteristics. More specifically, in the example illustrated in Table 1, when the selected memory cell **102** stores a logic 1 the sense amplifier **302** generates 251 UCLK signal pulses and 249 DCLK signal pulses each time T , and when the cell stores a logic 0 the sense amplifier generates 245 UCLK signal pulses and 255 DCLK signal pulses each time T . In Table 1, the far left column represents time intervals and the next column to the right represents the logic state of the bit stored in the selected memory cell **102** that is being sensed. The next column to the right represents the binary values of the SCNT count, and the far right column represents the decimal value of the corresponding binary SCNT count.

As illustrated in Table 1, the SCNT count is initially set to a binary value of 1000000000, which corresponds to a decimal value 512. After sensing the unknown bit stored in the selected memory cell **102** for the time T , which is a logic

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1 in the example of Table 1, the SCNT count has a binary value of 1000000010 or decimal value 514. After sensing the unknown bit for a second time T, the SCNT count has a binary value of 1000000100 and a decimal value of 516. At this point, the operation of the counter **306** is reversed as indicated in Table 1, which corresponds to the point at which the control circuit **316** activates the REV* signal, as previously described. A known logic 1 is stored in the selected memory cell **102** at this point and then sensed for the time T. After the time T, the SCNT count has the binary value of 1000000010 corresponding to 514 decimal. Thus, during the third sensing time T, the count of 2 added to the SCNT count during the second sensing time T is subtracted due to the reverse operation of the counter **306**. At this point, a known logic 0 is stored in the selected memory cell **102** and is sensed for the time T. After this time T, the SCNT count has the binary value 1000001100 which equals 524 decimal. At this point, the MSB bit of the SCNT count indicates the logic state of the unknown bit stored in the selected memory cell **102**, which equals a logic 1 in this example.

Another example of the operation of the read circuit **300** will now be described with reference to Table 2 set forth below.

TABLE 2

(Logic 1 = 251 UCLK, 249 DCLK; Logic 0 = 245 UCLK, 255 DCLK)			
Time	Sensed Bit	Counter Bits	Decimal Value
Start		1000000000	512
T	Unknown "0"	0111110110	502
T	Unknown "0"	0111101100	492
Switch UP and DN			
T	Known "1"	0111101010	490
T	Known "0"	0111110100	500

Table 2 illustrates the SCNT counts generated by counter **306** in the read circuit **300** after respective times T whether a logic 1 and 0 stored in the selected memory cell **102** have the specified characteristics. More specifically, in the example illustrated in Table 2, when the selected memory cell **102** stores a logic 1 the sense amplifier **302** generates 251 UCLK signal pulses and 249 DCLK signal pulses each time T, and when the cell stores a logic 0 the sense amplifier generates 245 UCLK signal pulses and 255 DCLK signal pulses each time T. In Table 2, the far left column represents time intervals and the next column to the right represents the logic state of the bit stored in the selected memory cell **102** that is being sensed. The next column to the right represents the binary values of the SCNT count, and the far right column represents the decimal value of the corresponding binary SCNT count.

As illustrated in Table 2, the SCNT count is initially set to a binary value of 1000000000, which corresponds to a decimal value 512. After sensing the unknown bit stored in the selected memory cell **102** for the time T, which is a logic 0 in the example of Table 2, the SCNT count has a binary value of 0111110110 or decimal value 502. After sensing the unknown bit for a second time T, the SCNT count has a binary value of 0111101100 and a decimal value of 492. At this point, the operation of the counter **306** is reversed as indicated in Table 2, which corresponds to the point at which the control circuit **316** activates the REV* signal, as previously described. A known logic 1 is stored in the selected memory cell **102** at this point and then sensed for the time T. After the time T, the SCNT count has the binary value of 0111101010 corresponding to 490 decimal. At this point, a

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known logic 0 is stored in the selected memory cell **102** and is sensed for the time T. After this time T, the SCNT count has the binary value 0111110100 which equals 500 decimal. At this point, the MSB bit of the SCNT count indicates the logic state of the unknown bit stored in the selected memory cell **102**, which equals a logic 0 in this example.

As illustrated by the examples of Tables 1 and 2, the MSB bit of the SCNT count equals a logic state of the unknown bit stored in the selected memory cell **102**. This can be understood by realizing that when the SCNT count is set to the initial value IV of 1000000000, the net of the UCLK and DCLK signal pulses applied to the counter **306** in both the normal and reverse modes of operation results in this initial value being incremented when the selected memory cell **102** stores a logic 1. As a result, the MSB bit of the SCNT count will remain a logic 1 in this situation. In contrast, when the SCNT count is set to the initial value IV of 1000000000, the net of the UCLK and DCLK signal pulses applied to the counter **306** in both the normal and reverse modes of operation results in this initial value being decremented when the selected memory cell **102** stores a logic 0. Accordingly, when the initial value IV is decremented, the MSB bit changes to a logic 0 and thus correctly indicates the data bit stored in the selected memory cell **102**.

The operation of the read circuit **300** may alternately be understood in the following terms. Assume that a logic 1 results in a net increase of the SCNT count by a value of X when sensed over a time T, and that a logic 0 results in a net change Y when sensed over the time T, where Y may be a net increase or decrease but where Y is less than X. In this situation, if the SCNT count has the initial value IV, then after 2T the SCNT count will have a value of either IV+2X or IV+2Y, depending on whether the sensed bit is a logic 1 or 0. If the values X and Y are then subtracted from the either IV+2X or IV+2Y, then a final value of the SCNT count will equal either a first final value FV1=IV+2X-X-Y or a second final value FV2=IV+2Y-X-Y. The first final value FV1=IV+X-Y, and since X>Y then FV1>IV. Thus, when the sensed bit is a logic 1, the SCNT count will have the value FV1, which is greater than IV. This means that if the initial value IV is selected with the MSB bit being a 1, then the MSB bit will equal a 1 when the SCNT count has the value FV1. Conversely, the second final value FV2=IV+Y-X, and since X>Y then FV2<IV. Thus, when the sensed bit is a logic 0, the SCNT count will have the value FV2, which is less than IV, meaning that if the initial value IV is selected with the MSB bit being a 1 then the MSB bit will be a 0 when the SCNT count has the value FV2.

Using these concepts to explain the specific operation of the read circuit **300** in the embodiment of FIG. 3, when a selected memory cell **102** stores a logic 1, the value of the SCNT count represents a logic 1 being sensed for the time 2T, which will equal a number greater than the initial value IV since a logic 1 includes more UCLK signal pulses than DCLK signal pulses. Accordingly, when the operation of the counter **306** is reversed, the known logic 1 subtracts $\frac{1}{2}$ the value of the SCNT count after the time 2T. When the known logic 0 is sensed, the reverse operation of the counter **306** results in the SCNT count after the time 3T being incremented. As a result, the SCNT count is greater than IV. Conversely, when a selected memory cell **102** stores a logic 0, the value of the SCNT count represents a logic 0 being sensed for the time 2T, which will equal a number less than the initial value IV since a logic 0 includes more DCLK signal pulses than UCLK signal pulses. Accordingly, when the operation of the counter **306** is reversed, the known logic 0 adds $\frac{1}{2}$ the value of the SCNT count after the time 2T.

When the known logic 1 is sensed, the reverse operation of the counter **306** results in the SCNT count after the time being decremented, leaving the final value of the SCNT count at a value less than the initial value IV and thus leaving the MSB bit equal to 0. One skilled in the art will realize that if the situation is reversed and a logic 1 includes more DCLK signal pulses than UCLK signal pulses relative to a logic 0, the read circuit **300** may operate in the identical manner except that the MSB bit of the SCNT count now represents the complement of the logic state of the data bit stored in a selected memory cell **102**.

The read circuit **300** generates the single MSB bit of the SCNT count to indicate the logic state of data stored in a selected memory cell **102**. With the read circuit **300**, there is no need to compare multiple N-bit count values to read data stored in a selected memory cell **102**. Moreover, with the read circuit **300**, there is no need to use a reference count RCNT, which will vary among selected memory cells **102**. Instead, the read circuit **300** utilizes actual count values for a known a logic 1 and logic 0 stored in the selected memory cell **102** in detecting a logic state of an unknown bit stored in the selected memory cell. This eliminates the need for a reference count that must be used with multiple selected memory cells **102** even though the precise counts corresponding to a logic 1 and a logic 0 vary among the memory cells, which introduces uncertainty in determining whether a given memory cell stores a logic 1 or a logic 0, as will be appreciated by those skilled in the art.

FIG. 5 is a simplified block diagram of a magnetic random access memory (MRAM) **500** including the read circuit **300** of FIG. 3 contained in a read/write circuit **502**. The read/write circuit **502** is coupled to a resistive memory-cell array, which is shown as the MRAM array **100** of FIG. 1, and operates during write operations to transfer data on a data bus DATA to addressed memory cells **102** in the MRAM array. During read operations, the read circuit **300** operates as previously described to read data from addressed memory cells in the MRAM array **100**, and the read/write circuit **502** places the read data onto the data bus DATA. The read/write circuit **502** also includes the circuitry for transferring a known logic 1 and 0 to selected memory cells **102** during read operations, as previously discussed with reference to the read circuit **300** of FIG. 3.

The MRAM **500** further includes an address decoder **504** that receives addresses from external circuitry (not shown), such as a processor or memory controller, on an address bus ADDR. In response to the received addresses, the address decoder **504** decodes the addresses and applies decoded address signals to access corresponding MRAM memory cells in the MRAM array **100**. A control circuit **506** applies a plurality of control signals **508** to control the MRAM array **100**, address decoder **504** and read/write circuit **502** during operation of the MRAM **500**. The control circuit **316** (FIG. 3) of the data read circuit **300** may be contained in the control circuit **506**, and thus the control signals **508** may include the REV*, EN, RST, and signals **318** discussed with reference to FIG. 3.

In operation, the external circuitry provides address, control, and data signals to the MRAM **500** over the respective ADDR, CONT, and DATA busses. During a write cycle, the external circuitry provides memory addresses on the ADDR bus, control signals on the CONT bus, and data on the DATA bus. In response to the control signals, the control circuit **506** generates controls signals **508**, including the REV signal, to control the MRAM array **100**, address decoder **504**, and read/write circuitry **504**. The address decoder **504** decodes the memory address on the ADDR bus and provides decoded

address signals to select the corresponding memory cells in the MRAM array **100**. The read/write circuitry **504** receives write data on the DATA bus, and applies the write data to the MRAM array **100** to store the data in the selected memory cells.

During a read cycle, the external circuitry provides a memory address on the ADDR bus and control signals on the CONT bus. Once again, in response to the control signals, the control circuit **506** generates controls signals **508** to control the MRAM array **100**, address decoder **504**, and read/write circuitry **504**. In response to the memory address, the address decoder **504** provides decoded address signals to access the corresponding memory cells in the array **100**. The read/write circuitry **504** provides data stored in the addressed memory cells onto the DATA bus to be read by the external circuit. One skilled in the art will understand circuitry for forming the address decoder **504**, read/write circuitry **504**, and control circuit **506**, and thus, for the sake of brevity, these components are not described in more detail.

Although only the single array **100** is shown in the MRAM **500**, the MRAM may include a plurality of arrays, and may also include additional components not illustrated in FIG. 5, as will be appreciated by those skilled in the art. Moreover, as previously mentioned, the MRAM **500** can include any type of resistive memory-cell array and is thus not limited to including MRAM arrays. In the MRAM **500**, the read circuit **300** is shown as included in the read/write circuit **502**, but one skilled in the art will realize portions of the read circuit could be included in the other components of the MRAM. For example, the control circuit **316** (FIG. 3) of the read circuit **300** could be included in the control circuit **506** of the MRAM **500**.

FIG. 6 is a block diagram of a computer system **600** including computer circuitry **602** that contains the MRAM **500** of FIG. 5. The computer circuitry **602** performs various computing functions, such as executing specific software to perform specific calculations or tasks. In addition, the computer system **600** includes one or more input devices **604**, such as a keyboard or a mouse, coupled to the computer circuitry **602** to allow an operator to interface with the computer system. Typically, the computer system **600** also includes one or more output devices **606** coupled to the computer circuitry **602**, such output devices typically being a printer or video display. One or more data storage devices **608** are also typically coupled to the computer circuitry **602** to store data or retrieve data from external storage media (not shown). Examples of typical storage devices **608** include hard and floppy disks, tape cassettes, compact disc read-only memories (CD-ROMs), read-write CD ROMs (CD-RW), and digital video discs (DVDs). Moreover, although the MRAM **500** is shown as being part of the computer circuitry **602**, the MRAM can also be used as a data storage device **608** since, as previously described, the nonvolatile nature and speed of the MRAM make it an attractive alternative to other storage media devices such as hard disks.

Even though various embodiments and advantages of the present invention have been set forth in the foregoing description, the above disclosure is illustrative only, and changes may be made in detail and yet remain within the broad principles of the invention. Therefore, the present invention is to be limited only by the appended claims.

The invention claimed is:

1. A method of sensing data stored in a selected resistive memory cell, the method comprising:

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generating a first count representing the logic state of an unknown initial data bit stored in the selected memory cell over a first time period equal to $(2 \times T)$;
 generating a second count representing a data bit having a first known logic state stored in the selected memory cell over a second time period equal to T ;
 generating a third count representing a data bit having a second known logic state stored in the selected memory cell over a third time period equal to T ; and
 determining the logic state of the unknown initial data bit stored in the selected memory cell from the first, second, and third counts.

2. The method of claim 1 wherein the operations of generating each comprise:

applying a current through the selected resistive memory cell and sensing the voltage developed responsive to the applied current;
 generating a plurality of up count signals and down count signals responsive to the sensed voltage; and
 adjusting the count responsive to the up count and down count signals.

3. The method of claim 1 wherein determining the logic state of the unknown initial data bit stored in the selected memory cell from the first, second, and third counts comprises subtracting the second and third counts from the first count to generate a final count.

4. The method of claim 3 further comprising determining the unknown logic state from a most significant bit of the final count.

5. The method of claim 1 wherein the resistive memory cell comprises an MRAM memory cell.

6. The method of claim 1 wherein the first known logic state comprises a logic 1 and the second known logic state comprises a logic 0.

7. The method of claim 6 wherein the second count representing a data bit having the first known logic state is greater than the third count representing a data bit having the second known logic state.

8. The method of claim 1 wherein generating a first count representing the logic state of an unknown initial data bit stored in the selected memory cell comprises generating an interim count representing the logic state of the unknown initial data bit stored in the selected memory cell over a time T , and thereafter doubling the interim count to generate the first count.

9. A method of sensing data stored in a selected resistive memory cell, the method comprising:

sensing an initial data bit stored in the selected memory over a first time period equal to $(2 \times T)$, the initial data bit having an unknown logic state;
 generating a count representing the logic state of the initial data bit responsive to the operation of sensing, the count having a most significant bit;
 transferring a data bit having a first known logic state into the selected memory cell;
 sensing the data bit having the known first logic state over a second time period equal to T ;
 adjusting the count responsive to sensing of the data bit having the first known logic state;
 transferring a data bit having a second known logic state into the selected memory cell;
 sensing the data bit having the second known logic state over a third time period equal to T ;
 adjusting the count responsive to sensing of the data bit having the second known logic state; and

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determining the unknown logic state of the initial data bit stored in the selected memory cell from the most significant bit of the count.

10. The method of claim 9 wherein the operations of sensing comprise applying a current through the selected resistive memory cell and sensing the voltage developed responsive to the applied current.

11. The method of claim 10 wherein the operations of sensing further comprise generating a plurality of up count signals and down count signals responsive to the sensed voltage.

12. The method of claim 11 wherein the operations of adjusting the count comprise incrementing the count responsive to each up count signal and decrementing the count responsive to each down count signal during the time $2T$.

13. The method of claim 12 wherein the operations of adjusting the count comprise decrementing the count responsive to each up count signal and incrementing the count responsive to each down count signal during sensing the bits having the first and second known logic states.

14. The method of claim 9 wherein determining the unknown logic state of the initial data bit stored in the selected memory cell from the most significant bit of the count comprises determining the unknown logic state equals the logic state of the most significant bit.

15. The method of claim 9 wherein the resistive memory cell comprises an MRAM memory cell.

16. The method of claim 9 wherein the first known logic state comprises a logic 1 and the second known logic state comprises a logic 0.

17. The method of claim 9 wherein generating a count representing the logic state of the initial data bit responsive to the operation of sensing includes setting the count to an initial value and adjusting the value of the count responsive to the operation of sensing.

18. A method of sensing data stored in a selected resistive memory cell, the method comprising:

setting a count value to an initial value;
 generating over a first time period equal to $(2 \times T)$ a first count representing the logic state of an unknown initial data bit stored in the selected memory cell;
 transferring a data bit having a first known logic state into the selected memory cell;
 generating over a second time period equal to T a second count representing the first known logic state of the data bit stored in the selected memory cell;
 transferring a data bit having a second known logic state into the selected memory cell;
 generating over a third time period equal to T a third count representing the second known logic state of the data bit stored in the selected memory cell;
 subtracting the second and third counts from the first count to obtain a final count; and
 determining the logic state of the unknown initial data bit stored in the selected memory cell from the final count.

19. The method of claim 18 wherein each of the operations of generating comprises:

applying a current through the selected resistive memory cell;
 sensing the voltage developed responsive to the applied current;
 generating a plurality of up count signals and down count signals responsive to the sensed voltage; and
 adjusting the corresponding count responsive to the up count and down count signals.

20. The method of claim 18 wherein determining the logic state of the unknown initial data bit stored in the selected

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memory cell from the final count comprises determining the unknown logic state from a most significant bit of the final count.

21. The method of claim 18 wherein the resistive memory cell comprises an MRAM memory cell.

22. The method of claim 18 wherein the first known logic state comprises a logic 1 and the second known logic state comprises a logic 0.

23. The method of claim 22 wherein the second count is greater than the third count.

24. The method of claim 18 wherein generating a first count representing the logic state of an unknown initial data bit stored in the selected memory cell comprises generating an interim count representing the logic state of the unknown initial data bit stored in the selected memory cell over a time T, and thereafter doubling the interim count to generate the first count.

25. A read circuit for sensing the logic state of a data bit stored in a resistive memory cell, the read circuit comprising:

a sense amplifier adapted to receive a sense voltage from a selected resistive memory cell, the sense amplifier operable to generate up and down clocking signals responsive to the sense voltage;

a counter including an up input and a down input, the counter operable to increment a count responsive to a clocking signal applied on the up input and operable to decrement the count responsive to a clocking signal applied on the down input; and

a switching circuit coupled between the sense amplifier and the counter and being adapted to receive a control signal, the switching circuit operable responsive to the control signal being inactive to apply the up and down clocking signals from the sense amplifier to the up and down inputs, respectively, of the counter, and operable responsive to the control signal being active to apply the up and down clocking signals from the sense amplifier to the down and up inputs, respectively, of the counter.

26. The read circuit of claim 25 wherein the switching circuit comprises:

a first pair of transmission gates coupled between the sense amplifier and the up and down inputs of the

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counter, the first pair of transmission gates applying the up and down clocking signals to the up and down inputs, respectively, when the control signal is inactive; and

a second pair of transmission gates coupled between the sense amplifier and the up and down inputs of the counter, the second pair of transmission gates applying the down and up clocking signals to the up and down inputs, respectively, when the control signal is active.

27. The read circuit of claim 25 further comprising a control circuit coupled to the sense amplifier, counter, and switching circuit, and operable to apply an enable signal to the sense amplifier to enable generation of the clocking signals by the sense amplifier, to apply a reset signal to the counter to reset the value of the count to an initial value, and to apply the control signal to the switching circuit, and operable to control the logic state of a data bit stored in the selected memory cell.

28. A read circuit for sensing the logic state of a data bit stored in a resistive memory cell, the read circuit comprising:

sensing means for sensing a sense voltage from a selected resistive memory cell and for generating up and down clocking signals responsive to the sense voltage;

counting means coupled to the sensing means for incrementing a count responsive to a first clocking signal and for decrementing the count responsive to a second clocking signal; and

switching means coupled to the sensing and counting means for applying the up and down clocking signals from the sensing means to the counting means as the first and second clocking signals, respectively, when a control signal is inactive, and for applying the down and up clocking signals from the sensing means to the counting means as the first and second clocking signals, respectively, when the control signal is active.

29. The read circuit of claim 28 wherein further comprising a control means coupled to the sensing, counting, and switching means for controlling the operation of the sensing, counting, and switching means and for controlling the logic state of data contained in the selected resistive memory cell.

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